## Performance

Will Crichton CS 242 – 11/26/18

### Key questions of performance

- 1. What programs need to be efficient?
- 2. How do we know if programs are efficient?
- 3. How can we make programs efficient?

### **Resource limits in 1975**



#### Apple II

- 32 KB of memory
- 1 MHz CPU
- 100 KB floppy disks
- \$5000+ (w/ inflation)

## The average programmer *always* has to care about performance.

### Extreme end: game development

Games were mostly engineered in a "data-oriented" way out of sheer necessity. There is not much room for abstraction when your target console has 128KB of RAM (SNES).

Every byte of storage is precious, so mostly games from this era are designed with very predictable manually managed layouts for their entire game state in memory. In the NES / SNES era, there was so little memory that generally the graphical representation of your game (tiles, sprites) and the logical representation of your game are the same,

In Mario 64, the entity structures are all exactly 608 bytes long, and there is a hard limit to 240 of them.

## **Resource limits today**

- 3 GHz processor
- 32+ GB of RAM
- GPU (1+ *tera*flop/s)
- 1 terabyte of disk

## The average programmer *rarely* has to care about performance.

## PL paradigms throughout history

#### • 1975: rise of systems languages

- Efficiency first: we don't have enough resources, carefully build a system

#### • 1995: rise of scripting languages

- Productivity first: systems languages are too hard (and we have the resources), so quickly glue together a system

#### • 2015: rise of functional languages

 Correctness first: scripting languages are too buggy, we need to know if our system works

### Latency: visual applications



Follow V

Replying to @rsnous @danluu

100%! I was part of a project trying to cut one frame of latency out of iOS's touch response. Cutting just a few ms down required changes top-to-bottom in the OS!

Alt story: iOS touch code had a fun rule – zero dynamic allocations allowed between touch HW event and app code!

3:53 PM - 31 Aug 2018



### Latency: data transfer



### Latency: media processing



**Figure 3:** Overview of our two processing pipelines. The input to both pipelines is a stream of Bayer mosaic (raw) images at full sensor resolution (for example, 12 Mpix) at up to 30 frames per second. When the camera app is launched, only the viewfinder (top row) is active. This pipeline converts raw images into low-resolution images for display on the mobile device's screen, possibly at a lower frame rate. In our current implementation the viewfinder is 1.6 Mpix and is updated at 15–30 frames per second. When the shutter is pressed, this pipeline suspends briefly, a burst of frames is captured at constant exposure, stored temporarily in main memory, and the software pipeline (bottom row) is activated. This pipeline aligns and merges the frames in the burst (sections 4 and 5), producing a single intermediate image of high bit depth, then applies color and tone mapping (section 6) to produce a single full-resolution 8-bit output photograph for compression and storage in flash memory. In our implementation this photograph is 12 Mpix and is computed in about 4 seconds on the mobile device.

#### Hasinoff et al. "Burst photography for high dynamic range and low-light imaging on mobile cameras". SIGGRAPH 2016

### Throughput: big data analytics















































### Huge perf gap for most programs



#### Saman Amarasinghe, MIT 6.172 "Performance Engineering". 2009

## **Throughput: simulation**



#### Resource usage: memory

| Process Name | Memory ~ |
|--------------|----------|
| Slack Helper | 365.8 MB |
| Slack Helper | 234.9 MB |
| 💁 Slack      | 91.8 MB  |
| Slack Helper | 89.8 MB  |

700 MB!

| Process Name         | $\textbf{Memory} \lor$ | Со |
|----------------------|------------------------|----|
| Google Chrome Helper | 738.8 MB               |    |
| Google Chrome Helper | 419.8 MB               |    |
| Google Chrome Helper | 399.4 MB               |    |
| Google Chrome Helper | 366.4 MB               |    |
| 🧿 Google Chrome      | 317.8 MB               |    |
| Google Chrome Helper | 57.2 MB                |    |
| Google Chrome Helper | 54.9 MB                |    |
| Google Chrome Helper | 53.5 MB                |    |
| Google Chrome Helper | 45.2 MB                |    |
| Google Chrome Helper | 44.5 MB                |    |
| Google Chrome Helper | 44.0 MB                |    |
| Google Chrome Helper | 43.8 MB                |    |
| Google Chrome Helper | 43.6 MB                |    |
| Google Chrome Helper | 43.5 MB                |    |
| Google Chrome Helper | 43.3 MB                |    |
| Google Chrome Helper | 43.2 MB                |    |
| Google Chrome Helper | 43.1 MB                |    |
| Google Chrome Helper | 42.8 MB                |    |
| Google Chrome Helper | 40.0 MB                |    |
| Google Chrome Helper | 30.1 MB                |    |
| Google Chrome Helper | 18.8 MB                |    |
| Google Chrome Helper | 18.4 MB                |    |
| Google Chrome Helper | 17.3 MB                |    |

#### 3 GB!!!

## Performance profiling is a dark art

- Perf, cProfile, etc. are rarely taught in school
- Debuggers are great, but a lot of profiling can be pretty close to "printf debugging"
- A lot more work to do here!
- (Claim: "Profiling", "Debugging" should be required courses in any CS curriculum.)

#### Linux Performance Tools



## How to improve program performance:

Change the program
 Change the programmer

#### Automatic optimization

- Promise of automatic optimizers: don't worry about performance, we'll take care of everything
- Type checks, memory management, parallelization
- Issue is <u>consistency</u>

## Algebraic simplification

- Some statements can be deleted
   x := x + 0
   x := x \* 1
- Some statements can be simplified

## **Register allocation**



## Free RegistersR\_0R\_1R\_2

Stanford CS 143 "Compilers"

## Hot loop compilation

- Runtime profiling to identify hot spots in code
- Execute compiler during runtime
- Swap in compiled code for runtime code
- Monomorphization for dynamically typed code

#### Virtual method call optimization

invokevirtual

in Java



Stanford CS 242, 2016

#### Virtual method call optimization

invokevirtual





I trawled V8 issue tracker and found few

У,

- Issue 6391: StringCharCodeAt slow
- Issue 7092: High overhead of Strip
- Issue 7326: Performance degrad

Overhead Symbo

Javasc

\*compareByOriginalPositions ../dist/source-map Builtin:ArgumentsAdaptor na 4.49% Builtin:CallFunction\_ReceiverIsNullOrUndefined 17.02% \*compareByGeneratedPositionsDeflated ../dist/source-map.js:1063 11.20% \*SourceMapConsumer\_parseMappings ../dist/source-map.js:1894 \*SourceMapConsumer\_parseMappings ../dist/source-map.js:1894 3.58% 2.11% Builtin:StringEqual 1.25% v8::internal::StringTable::LookupStringIfExists\_NoAllocate \*SourceMapConsumer\_parseMappings ../dist/source-map.js:1 Builtin:Call\_ReceiverIsNullOrUndefined Builtin:StringCharAt v8::internal::(anonymous namespace)::StringTableNoAl ŀ, 1.22% 1.21% 1.16% Builtin:StringPrototypeSlice v8::internal::(anonymous namespace)::MakeStringThin Builtin:KeyedLoadIC\_Megamorphic 1.14% v8::internal::(anonymous namespace)::CopyObjectToObjectElements 0.90% v8::internal::String::VisitFlat<v8::internal::IteratingStringHasher> 0.86% 0.82% \*SourceMapConsumer\_parseMappings ../dist/source-map.js:1894 0.80% 0.76% \*doQuickSort ../dist/source-map.js:2752 v8::internal::IncrementalMarking::RecordWriteSlow 0.72% 0.68% 0.64% 0.56%



#### Scalar

| #define T double                               |
|--|
| <pre>void add(T* x, T* y, T* z, int N) {</pre> |
| <pre>for(int i = 0; i &lt; N; ++i) {</pre>     |
| T x1, y1, z1;                                  |
| x1 = x[i];                                     |
| y1 = y[i];                                     |
| z1 = x1 + y1;                                  |
| z[i] = z1;                                     |
| }  |
| }  |



#define T double
void add(T\* x, T\* y, T\* z, int N) {
for(int i = 0; i < N; i += 4) {
 \_\_m256d x1, y1, z1;
 x1 = \_mm256\_loadu\_pd(x + i);
 y1 = \_mm256\_loadu\_pd(y + i);
 z1 = \_mm256\_add\_pd(x1, y1);
 \_mm256\_storeu\_pd(z + i, z1);
}</pre>

| l |  |  |
|---|--|--|
|   |  |  |



#### **Auto-vectorization**

Steeped in generating great code for regular loops that performed dense matrix math, for a long time most of the Intel compiler team denied that anything more than their auto-vectorizer was needed to take care of vector unit utilization. We quickly fell into a cycle:

- They'd inform the graphics folks that they'd improved their auto-vectorizer in response to our requests and that it did everything we had asked for.
- We'd try it and find that though it was better, boy was it easy to write code that wasn't actually compiled to vector code—it'd fail unpredictably.
- We'd give them failing cases, a few months would would pass and they'd inform us that the latest version solved the problem.

And so on.

It didn't take much to fall off the vectorization path. They tried to patch things up at first but eventually they threw up their hands and came up with #pragma simd, which would disable the "is it safe to vectorize this" checks in the auto-vectorizer and vectorize the following loop no matter what. (Once a #pragma is proposed to solve a hard problem, you know things aren't in a good place.)

#### Play with the compiler flags

- ≻icc –help
- Find the best flags
  - icc -c -O3 -xT -msse3 mxm.c
- Use information from icc
  - icc -vec-report5 ...
- Generate assembly and stare!
  - Icc -S -fsource-asm -fverboseasm...

#### Tweaked the program until the compiler is happy 🛞

```
;;; for (j2 = 0; j2 < N; j2 += BLOCK X)
        xorl
                   %edx, %edx
        xorl
                   %eax, %eax
                   %xmm0, %xmm0
        xorps
;;; for (k2 = 0; k2 < N; k2 += BLOCK Y)
;;; for(i = 0; i < N; i++)
        xorl
                   %ebx, %ebx
        xorl
                   %ecx, %ecx
;;; for(j = 0; j < BLOCK X; j++)
                   %r9d, %r9d
        xorl
;;; for (k = 0; k < BLOCK Y; k++)
;;; IND (A,i,j+j2,N) += IND (B,i,k+k2,N) * IND (Cx,j+j2,k+k2,N);
                   %ecx, %r8
        movslq
                   (%rdx,%rcx), %esi
        lea
        movslq
                   %esi, %rdi
                   $3, %rdi
        shlq
                   %eax, %rsi
        movslq
        shlq
                   $3, %rsi
..B1.13:
                   %xmm0, %xmm2
        movaps
                   A(%rdi), %xmm1
        movsd
        xorl
                   %r10d, %r10d
..B1.14.
        movaps
                   B(%r10,%r8,8), %xmm3
        mulpd
                   Cx(%r10,%rsi), %xmm3
        addpd
                   %xmm3, %xmm1
                                                      instructions
                   16+B(%r10,%r8,8), %xmm4
        movaps
        mulpd
                   16+Cx(%r10,%rsi), %xmm4
        addpd
                   %xmm4, %xmm2
        movaps
                   32+B(%r10,%r8,8), %xmm5
        mulpd
                   32+Cx(%r10,%rsi), %xmm5
        addpd
                   %xmm5, %xmm1
        movaps
                   48+B(%r10,%r8,8), %xmm6
        mulpd
                   48+Cx(%r10,%rsi), %xmm6
        addpd
                   %xmm6, %xmm2
        movaps
                   64+B(%r10,%r8,8), %xmm7
                                                     Ш
                   64+Cx(%r10,%rsi), %xmm7
        mulpd
                                                     ົ
        addpd
                   %xmm7, %xmm1
                                                     S
        movaps
                   80+B(%r10,%r8,8), %xmm8
                   80+Cx(%r10,%rsi), %xmm8
        mulpd
                                                     Inner loop:
        addpd
                   %xmm8, %xmm2
        movaps
                   96+B(%r10,%r8,8), %xmm9
        mulpd
                   96+Cx(%r10,%rsi), %xmm9
        addpd
                   %xmm9, %xmm1
        movaps
                   112+B(%r10,%r8,8), %xmm10
        mulpd
                   112+Cx(%r10,%rsi), %xmm10
        addpd
                   %xmm10, %xmm2
        addq
                   $128, %r10
                   $8192, %r10
        cmpq
                   ..B1.14
        jl
                                  # Prob 99%
```

# Auto-vectorization is not a programming model!

## Thesis:

## The future of performance optimization is better programming models, not better optimizers.

## Parallel processing of large datasets with Spark

#### Word count

## Parallel processing of large datasets with Spark

#### **K-means clustering**

```
lines = spark.read.text(sys.argv[1]).rdd.map(lambda r: r[0])
data = lines.map(parseVector).cache()
K = int(sys.argv[2])
convergeDist = float(sys.argv[3])
kPoints = data.takeSample(False, K, 1)
tempDist = 1.0
while tempDist > convergeDist:
    closest = data.map(
        lambda p: (closestPoint(p, kPoints), (p, 1)))
    pointStats = closest.reduceByKey(
        lambda p1_c1, p2_c2: (p1_c1[0] + p2_c2[0], p1_c1[1] + p2_c2[1]))
    newPoints = pointStats.map(
        lambda st: (st[0], st[1][0] / st[1][1])).collect()
    tempDist = sum(np.sum((kPoints[iK] - p) ** 2) for (iK, p) in newPoints)
    for (iK, p) in newPoints:
        kPoints[iK] = p
```

## Low-level optimization of image processing with Halide



machinelearninguru.com. "Image Filtering"

## Low-level optimization of image processing with Halide

Func blur\_3x3(Func input) {
 Func blur\_x, blur\_y;
 Var x, y, xi, yi;

// The algorithm - no storage or order blur\_x(x, y) = (input(x-1, y) + input(x, y) + input(x+1, y))/3; blur\_y(x, y) = (blur\_x(x, y-1) + blur\_x(x, y) + blur\_x(x, y+1))/3;

#### return blur\_y;

## Low-level optimization of image processing with Halide

Fredo Durand, "High-Performance Image Processing"

#### Halide

0.9 ms/megapixel

Func box\_filter\_3x3(Func in) {
 Func blurx, blury;
 Var x, y, xi, yi;

// The algorithm - no storage, order blurx(x, y) = (in(x-1, y) + in(x, y) + in(x+1, y))/3; blury(x, y) = (blurx(x, y-1) + blurx(x, y) + blurx(x, y+1))/3;

// The schedule - defines order, locality; implies storage
blury.tile(x, y, xi, yi, 256, 32)
 .vectorize(xi, 8).parallel(y);
blurx.compute\_at(blury, x).store\_at(blury, x).vectorize(x, 8);

return blury;

#### C++

#### 0.9 ms/megapixel

```
void box_filter_3x3(const Image &in, Image &blury) {
   m128i one third = mm set1 epi16(21846);
  #pragma omp parallel for
  for (int yTile = 0; yTile < in.height(); yTile += 32) {</pre>
     m128i a, b, c, sum, avg;
      _m128i blurx[(256/8)*(32+2)]; // allocate tile blurx array
    for (int xTile = 0; xTile < in.width(); xTile += 256) {</pre>
      __m128i *blurxPtr = blurx;
      for (int y = -1; y < 32+1; y++) {</pre>
        const uint16 t *inPtr = &(in[yTile+y][xTile]);
        for (int x = 0; x < 256; x += 8) {
         a = _mm_loadu_si128((__m128i*)(inPtr-1));
         b = mm loadu si128((__m128i*)(inPtr+1));
         c = _mm_load_si128((__m128i*)(inPtr));
         sum = _mm_add_epi16(_mm_add_epi16(a, b), c);
         avg = _mm_mulhi_epi16(sum, one_third);
         _mm_store_si128(blurxPtr++, avg);
         inPtr += 8;
      }}
      blurxPtr = blurx;
      for (int y = 0; y < 32; y++) {</pre>
          m128i *outPtr = (__m128i *)(&(blury[yTile+y][xTile]));
        for (int x = 0; x < 256; x += 8) {
          a = _mm_load_si128(blurxPtr+(2*256)/8);
          b = mm load_si128(blurxPtr+256/8);
          c = mm load si128(blurxPtr++);
          sum = _mm_add_epi16(_mm_add_epi16(a, b), c);
          avg = _mm_mulhi_epi16(sum, one_third);
          _mm_store_si128(outPtr++, avg);
```

#### Low-level optimization of image processing with Halide

**Reference: 300 lines C++** 

Adobe: 1500 lines 3 months of work 10x faster (vs. reference)

Halide: 60 lines 1 intern-day

20x faster (vs. reference) **2x faster** (vs. Adobe)

GPU: 70x faster (vs. reference)

Fredo Durand, "High-Performance Image Processing"



Histogram **v** 

+24

### **Efficient DOM updates with React**

```
var toggled = false;
$('button').click(function() {
                                            $('button').click(function() {
  if (!$('#elem1').is(':visible')) {
                                              toggled = !toggled;
    $('#elem1').show();
                                              render();
    $('#elem2').show();
                                            });
  } else {
    $('#elem1').hide();
                                            function render() {
   $('#elem2').hide();
                                             let html = toggled
  }
                                               ? '<div id="elem1">Elem 1</div>' +
});
                                                  '<div id="elem2">Elem 2</div>'
                                                : 11:
                                              $('#container').html(html);
                                            }
   class Container extends React.Component {
      state = {toggled: false}
      render() {
        return <div id="container">
          <button onClick={</pre>
            () => this.setState({toggled: !this.state.toggled})}>
            Click me
          </button>
          {this.state.toggled
           ? <div><div>Elem 1</div><div>Elem 2</div>
           : <div />}
        </div>;
   ReactDOM.render(<Container />, document.getElementById('container'));
```

#### Domain models aren't a panacea

The published work on big data systems has fetishized scalability as the most important feature of a distributed data processing platform... Contrary to the common wisdom that effective scaling is evidence of solid systems building, any system can scale arbitrarily well with a sufficient lack of care in its implementation.

We offer a new metric for big data platforms, COST, or the Configuration that Outperforms a Single Thread. The COST of a given platform for a given problem is the hardware configuration required before the platform outperforms a competent singlethreaded implementation.

#### Domain models aren't a panacea

```
fn PageRank20(graph: GraphIterator, alpha: f32) {
    let mut a = vec![0f32; graph.nodes()];
    let mut b = vec![0f32; graph.nodes()];
    let mut d = vec![0f32; graph.nodes()];
    graph.map_edges(|x, y| { d[x] += 1; });
    for iter in 0..20 {
      for i in 0..graph.nodes() {
         b[i] = alpha * a[i] / d[i];
         a[i] = 1f32 - alpha;
      }
      graph.map_edges(|x, y| { a[y] += b[x]; });
    }
}
```

}

#### ← 16 lines of Rust

| scalable system     | cores | twitter | uk-2007-05 |
|---------------------|-------|---------|------------|
| GraphChi [12]       | 2     | 3160s   | 6972s      |
| Stratosphere [8]    | 16    | 2250s   | _          |
| X-Stream [21]       | 16    | 1488s   | -          |
| Spark [10]          | 128   | 857s    | 1759s      |
| Giraph [10]         | 128   | 596s    | 1235s      |
| GraphLab [10]       | 128   | 249s    | 833s       |
| GraphX [10]         | 128   | 419s    | 462s       |
| Single thread (SSD) | 1     | 300s    | 651s       |
| Single thread (RAM) | 1     | 275s    | -          |



- Despite hardware improvements, performance still matters
- Changing the programmer > changing the program
- General optimizations are hard, domain optimizations are easier